# 6. Probabilities: Markov chains and statistical model checking

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System Verification (CC4084) 2024/2025

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https://fm-dcc.github.io/sv2425



Where we are

# **Syllabus**



- Introduction to model-checking
- CCS: a simple language for concurrency
  - Syntax
  - Semantics
  - Equivalence
  - mCRL2: modelling
- Dynamic logic
  - Syntax
  - Semantics
  - Relation with equivalence
  - mCRL2: verification

- Timed Automata
  - Syntax
  - Semantics (composition, Zeno)
  - Equivalence
  - UPPAAL: modelling
- Temporal logics (LTL/CTL)
  - Syntax
  - Semantics
  - UPPAAL: verification
- Probabilistic and stochastic systems
  - Going probabilistic
  - UPPAAL: monte-carlo

# **Going probabilistic**

# **Motivation**



#### Systems can get very complex

- E.g., 5 components, 3 possible traces each
- No communication (pure interleaving)
- Many permutations

# **Motivation**



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# **Motivation**



### Systems can get very complex

- E.g., 5 components, 3 possible traces each
- No communication (pure interleaving)
- Many permutations
- More components, more traces untreatable
- Verifying deadlock freedom (and others) requires traversing all states
- Approximation:
  - traverse only part of the states
  - give more priority to some actions
  - return (statistically) likelihood of a given property

FC

- $\alpha: S \rightarrow N \times S$  Moore machine
- $\alpha: S \to \text{Bool} \times S^N$  deterministic automata
- $\alpha: S \to \text{Bool} \times P(S)^N$  non-deterministic automata (reactive)
- $\alpha: S \to P(N \times S)$  non deterministic LTS (generative)
- $\alpha: S \rightarrow (S+1)^N$  partial deterministic LTS
- $\alpha: S \to P(S)$  unlabelled TS
- $\alpha: S \to D(S)$  Markov chain



#### **Markov chains**

$$\alpha: S \to \mathrm{D}(S)$$

where D(S) is the set of all discrete probability distributions on set S

A Markov chain goes from a state s to a state s' with probability p if

$$lpha(s)=\mu$$
 with  $\mu(s')=p>0$ 



#### Recall

 $\mu: \mathcal{S} \rightarrow [0,1]$  is a discrete probability distribution if

- $\{s \in S \mid \mu(s) > 0\}$ , is finite (called the support of  $\mu$ ), and
- $\sum_{s\in S}\mu(s)=1$

# **Examples**

- Dirac distribution:  $\mu_s^1 = \{s \to 1\}$
- Product distribution:  $(\mu_1 imes \mu_2) \langle s, t 
  angle = \mu_1(s) imes \mu_2(t)$

Example





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**Reactive PTS** 



$$\alpha: S \to (\mathrm{D}(S) + 1)^N$$

### Ex. 6.1: Formalise the systems below as functions



#### Notions of bisimulation arise naturally.

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**Generative PTS** 



$$\alpha: S \to \mathrm{D}((S \times N) + 1)$$



Ex. 6.2: Now (generative) – formalise it a[0.4] a[0.4] b[0.3]b[0.2]

 $A_2$ 

# **Probabilities in Uppaal**

# **Stochastic Timed Automata – examples**







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Probabilities in Uppaal



# $\langle \textbf{L}, \textbf{L}_0, \textbf{Act}, \textbf{C}, \textbf{Tr}, \textbf{Inv} \rangle$

where

- L is a set of locations, and  $L_0 \subseteq L$  the set of initial locations
- Act is a set of actions and C a set of clocks
- $Tr \subseteq L \times C(C) \times Act \times P(C) \times \mathbb{N} \times L$  is the transition relation

$$\ell_1 \xrightarrow{g,a,U,w} \ell_2$$

denotes a transition from location  $\ell_1$  to  $\ell_2$ , labelled by *a*, enabled if guard *g* is valid, which, when performed, resets the set *U* of clocks, with a probability given by the weight *w* 

Inv : L → C(C) + Q is the assignment of invariants or rates (of an exponential distribution) to locations

where C(C) denotes the set of clock constraints over a set C of clock variables

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Probabilities in Uppaal







- Probability of  $\langle A1_0, \overline{0} \rangle \xrightarrow{0.5} \langle A1_0, \overline{0.5} \rangle$ ?
- Probability of  $\langle A2_0, \overline{0} \rangle \xrightarrow{0.5} \langle A2_0, \overline{0.5} \rangle$ ?







- Probability of  $\langle A1_0, \overline{0} \rangle \xrightarrow{0.5} \langle A1_0, \overline{0.5} \rangle$ ?
- Probability of  $\langle A2_0, \overline{0} \rangle \xrightarrow{0.5} \langle A2_0, \overline{0.5} \rangle$ ?
- Probability of  $\langle A3_0, \overline{0} \rangle \xrightarrow{0.5} \langle A3_0, \overline{0.5} \rangle$ ?

# Again A1,A2,A3: Timed PTS







- Probability of  $\langle A1_0, \overline{0} \rangle \xrightarrow{0.5} \langle A1_0, \overline{0.5} \rangle$ ?
- Probability of  $\langle A2_0, \overline{0} \rangle \xrightarrow{0.5} \langle A2_0, \overline{0.5} \rangle$ ?
- Probability of  $\langle A3_0, \overline{0} \rangle \xrightarrow{0.5} \langle A3_0, \overline{0.5} \rangle$ ?
- Probability of reaching A1<sub>1</sub>?
- Probability of reaching A2<sub>1</sub>?

# Again A1,A2,A3: Timed PTS







- Probability of  $\langle A1_0, \overline{0} \rangle \xrightarrow{0.5} \langle A1_0, \overline{0.5} \rangle$ ?
- Probability of  $\langle A2_0, \overline{0} \rangle \xrightarrow{0.5} \langle A2_0, \overline{0.5} \rangle$ ?
- Probability of  $\langle A3_0, \overline{0} \rangle \xrightarrow{0.5} \langle A3_0, \overline{0.5} \rangle$ ?
- Probability of reaching A1<sub>1</sub>?
- Probability of reaching A2<sub>1</sub>?

= ....

Probability of reaching A3<sub>END</sub> in less than 4.3?

# A1: When does it end?





- Run 102000 times
- Histogram: how many times it took [9..9.1] seconds?
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# A2: When does it end?







- Run 100000 times
- Histogram: how many times it took [9..9.1] seconds?

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# A3: When does it end?







- Run 300000 times
- Histogram: how many times it took [9..9.1] seconds?
  - ...

# **Generative Timed PTS**



$$\begin{array}{c} \alpha: \mathcal{S} \to \mathrm{D}_{\textit{disc}}((\mathcal{S} \times \mathcal{N}) + 1) \\ \alpha: \mathcal{S} \to \mathrm{D}_{\textit{cont}}((\mathcal{S} \times (\mathcal{N} + \mathcal{R}_0^+) + 1)) \end{array}$$



Ex. 6.3: Now (Timed PTS) – formalise it



# **Generative Timed PTS**



$$\begin{array}{l} \alpha: S \to \mathrm{D}_{disc}((S \times N) + 1) \\ \alpha: S \to \mathrm{D}_{cont}((S \times (N + \mathcal{R}_0^+) + 1)) \end{array}$$

#### Notes

- Continuous time: continuous distribution
- Probabilities both at continuous delays and discrete transitions.

# Ex. 6.3: Now (Timed PTS) – formalise it



Probabilistic queries in Uppaal



- Pr[c<=10; 100] ([] safe) runs 100 stochastic simulations and estimates the probability of safe remaining true within 10 cost units, based on 100 runs.</li>
- Pr[<=10] (<> good) runs a number of stochastic simulations and estimates the probability of good eventually becoming true within 10 time units. The number of runs is decided based on the probability interval precision (±ε) and confidence level (level of significance α).
- Pr[<=10] (<> good) >= 0.5 checks if the probability of reaching good within 10 time units is greater than 50% (less runs than calculating the probability, using "Walds's algorithm")
- E[<=10; 100] (max: cost) runs 100 stochastic simulations and estimates the maximal value of cost expression over 10 time units of stochastic simulation.

More at https://docs.uppaal.org/language-reference/query-syntax/statistical\_queries/

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Probabilistic queries in Uppaal

# Running a single simulation



- simulate [<=10] { x, y } creates one stochastic simulation run of up to 10 time units in length and plot the values of x and y expressions over time (after checking, right-click the query and choose a plot).</li>
- Variations: [c<=10] / [#<=10] based on clock c or based on the number of transitions.

# **Replicate the histograms**







### Ex. 6.4: Replicate the visualisation

# $\ensuremath{\text{Ex. 6.5:}}\xspace$ Replicate the visualisation also for A1 and A3

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**Ex. 6.6:** Adapt the model to make it stochastic

# Ex. 6.7: Adapt requirements to make them probabilistic

- 1. The lamp can become bright;
- 2. The lamp will eventually become bright;
- 3. The lamp can never be on for more than 3600s;
- 4. It is possible to never turn on the lamp;
- 5. Whenever the light is bright, the clock *y* is non-zero;
- 6. Whenever the light is bright, it will eventually become off.