3. Concurrency in Java and its memory model

Nelma Moreira & José Proença Concurrent programming (CC3040) 2024/2025

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https://fm-dcc.github.io/cp2425







Overview



Blocks of sequential code running concurrently and sharing memory:

- What is Scala?
- Concurrency in Java and its memory model
- Basic concurrency blocks and libraries
- Futures and promises
- Actor model



Synchronisation

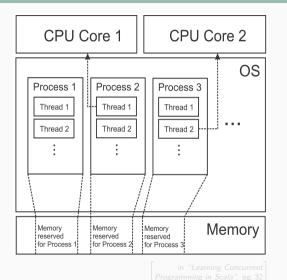
- Coordination of multiple executions in
- a concurrent system
- Mechanisms to order concurrent executions
- Mechanisms to exchange information

Exchanging information

- Concurrent programs: shared memory communication
- Distributed programs: message passing communication

Processes and threads





Starting a new JVM instance always creates only one process.

In that process, multiple threads can run simultaneously.

Unlike runtimes (e.g. Python), the JVM: does not implement its custom threads, maps each Java thread to an OS thread

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Managing threads



```
object ThreadsMain extends App {
  val t: Thread =
    Thread.currentThread
  val name = t.getName
  println(s"I_am_the_thread_$name")
}
```

Using SBT, this prints: [info] I am the thread sbt-bg-threads-1

```
In SBT do "set fork := true"
(or add "fork := true" to build.sbt)
It will then it prints:
[info] I am the thread main
```

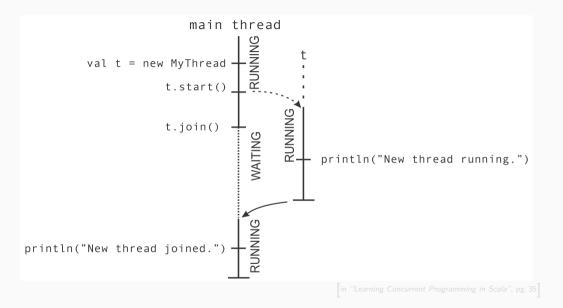


```
object ThreadsCreation extends App {
   class MyThread extends Thread {
      override def run(): Unit = {
        println("New_uthread_urunning.")
      }
   }
   val t = new MyThread
   t.start()
   t.join()
   println("New_uthread_ujoined.")
}
```

start eventually causes
run to execute in a new thread;

the OS decides when;

join puts the main thread in a waiting state, and allows the OS to re-assign the processor.





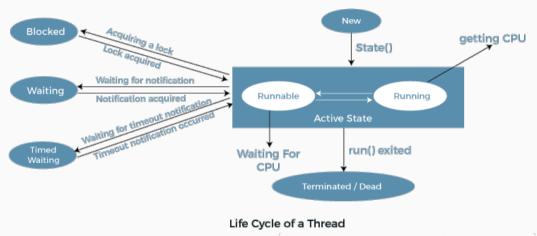
```
def thread(body: =>Unit): Thread = {
  val t = new Thread {
    override def run() = body
  }
  t.start()
  t
}
```



```
def thread(body: =>Unit): Thread = {
  val t = new Thread {
    override def run() = body
  }
  t.start()
  t
}
```

Using the thread function

```
object ThreadsSleep extends App {
  val t = thread {
    Thread. sleep(1000)
    log("New_thread_running.")
    Thread. sleep(1000)
    log("Still_running.")
    Thread. sleep(1000)
    log("Completed.")
  }
 t.join()
  log("New__thread__joined.")
```



n https://static.javatpoint.com/core/images/life-cycle-of-a-thread.png



What to expect?



```
object ThreadsNondeterminism
    extends App {
    val t = thread {
        log("New_thread_running.")
    }
    log("...")
    log("...")
    t.join()
    log("New_thread_joined.")
}
```

What to expect?

- "New thread joined" printed always at the end
- Other prints not always in the same order – nondeterministic execution



```
object ThreadsNondeterminism
    extends App {
    val t = thread {
        log("New_thread_running.")
    }
    log("...")
    log("...")
    t.join()
    log("New_thread_joined.")
}
```

What to expect?

- "New thread joined" printed always at the end
- Other prints not always in the same order – nondeterministic execution
- Common in concurrent applications what makes it so hard
- Note: join also forces all memory writes from the threads before proceeding

Control of the execution order



action α happens-before (HB) action β means action β sees the memory writes of action α

- **Program order:** α in a thread HB every subsequent β in that program and thread
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- Transitivity: if $\alpha \text{ HB } \beta$ and $\beta \text{ HB } \gamma$, then $\alpha \text{ HB } \gamma$

•



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- · · · ·

Data race: when a write to memory does not happen-before its intended read.



- join provides guarantees that other threads terminated
- Not enough we may want to inform other treads without terminating

Example 1: shared counter for unique IDs

```
object ThreadsUnprotectedUid extends App {
  var uidCount = OL
  def getUniqueId() = {
    val freshUid = uidCount + 1
    uidCount = freshUid
    freshUid
  }
...
}
```

What can go wrong?

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```
...
def printUniqueIds(n: Int): Unit = {
  val uids = for (i<- 0 until n)
      yield getUniqueId()
  log(s"Generated_uids:_u$uids")
}
val t = thread { printUniqueIds(5) }
printUniqueIds(5)
t.join()
...</pre>
```

What do you expect?

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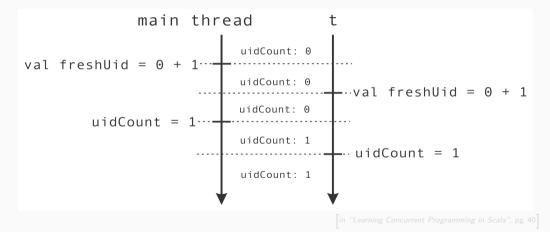
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printUniqueIds(5)
t.join()
...</pre>
```

Race Condition

when the output of a concurrent program depends on how the statements are scheduled.







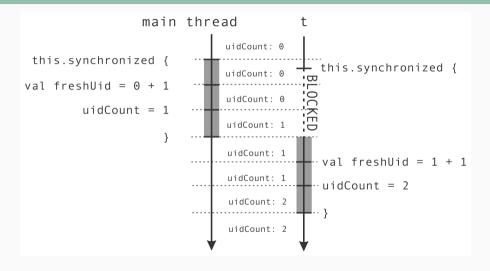


synchronized is:

- a fundamental Scala/Java construct for atomic executions
- can be called in any object (or instance of a class)
- ensures atomic execution wrt the object
- We say obj.synchronized
 - acquires the lock/monitor of obj at the start
 - releases the lock/monitor of obj at the end

```
def getUniqueId() =
   this.synchronized {
      val freshUid = uidCount + 1
      uidCount = freshUid
      freshUid
   }
```

Updating counter in parallel atomically

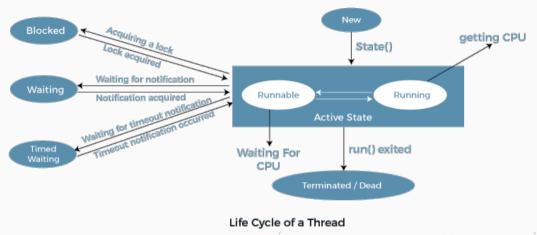


in "Learning Concurrent Programming in Scala", pg. 41

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Control of the execution order





n https://static.javatpoint.com/core/images/life-cycle-of-a-thread.png



- using the synchronized statement has some (not too large) overhead
- not using synchronized can easily lead to errors, even if all seems correct

Find the bug in the next slide...

Find the bug



```
object ThreadSharedStateAccessReordering extends App {
 for (i <- 0 until 100000) {</pre>
    var a = false
    var b = false
   var x = -1
   var y = -1
    val t1 = thread {
      a = true
      y = if (b) 0 else 1
    3
    val t2 = thread {
      b = true
      x = if (a) 0 else 1
    }
   t1.join()
   t2.join()
    assert(!(x==1 && y==1), s"x=$x, y=$y")
  }
}
```



- The previous code can raise an error: both x and y can become 1!
- JVM can reorder statements in a thread when they seem to be independent.
- Because some processors do not always execute instructions in the expected order, to increase performance.
- (Known as "weak memory model")
- A synchronized block would solve this:
 - also enclosing each assignment in a synchronized block
 - synchronized sets up a memory barrier



- every object has a *lock*
- a running thread can aquire multiple locks from different objects

Example 2: Logging Bank Transfers

```
object SynchronizedNesting extends App {
   import scala.collection.
   private val transfers = mutable.ArrayBuffer[String]()
   def logTransfer(name: String, n: Int) = transfers.synchronized {
     transfers += s"transfer_to_account_t'sname'_=:sn"
   }
   class Account(val name: String, var monev: Int)
   def add(account: Account, n: Int) = account.synchronized {
       account.money += n
       if (n > 10) logTransfer(account.name, n)
   }
   . . .
```



```
private val transfers = mutable.ArrayBuffer[String]()
def logTransfer(name: String, n: Int) = transfers.synchronized {
 transfers += s"transfer..to..account...'$name'..=..$n"
3
class Account(val name: String, var money: Int)
def add(account: Account, n: Int) = account.synchronized {
    account.money += n
   if (n > 10) logTransfer(account.name, n)
}
val jane = new Account("Jane", 100)
val john = new Account("John", 200)
val t1 = thread { add(jane, 5) }
val t2 = thread { add(john, 50) }
val t3 = thread { add(jane, 70) } // will not corrupt Jane's account
t1.join(); t2.join(); t3.join()
log(s"---__transfers___---\n$transfers")
```



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- Transitivity: if α HB β and β HB γ , then α HB γ
- · · · ·
- Monitor locking: unlocking HB every subsequent locking (of the same lock)

Deadlocks



Deadlock

when two or more executions wait for each other before proceeding

- Studied in the first module with prof. Nelma Moreira
- Often caused by locks that are not released at the right time

```
object SynchronizedDeadlock extends App {
  import SynchronizedNesting.Account
  def send(a: Account, b: Account, n: Int) = a.synchronized {
     b.synchronized {
        a.money -= n
        b.money += n
     }
  }
  ... // can this go wrong?
}
```



```
def send(a: Account, b: Account, n: Int) = a.synchronized {
  b.synchronized {
    a.money -= n
    b.monev += n
 }
val l = new Account("Lucy", 1000)
val j = new Account("Jim", 2000)
val t1 = thread { for (i<- 0 until 100) send(1, j, 1) }
val t2 = thread { for (i<- 0 until 100) send(j, 1, 1) }</pre>
t1. join(); t2. join()
log(s"a_{\sqcup}=_{\sqcup} {a.money}, _{\sqcup}b_{\sqcup}=_{\sqcup} {b.money}")
```



```
def send(a: Account, b: Account, n: Int) = a.synchronized {
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It works but...



```
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val t2 = thread { for (i<- 0 until 100) send(j, 1, 1) }</pre>
t1. join(); t2. join()
log(s"a_{\sqcup}=_{\sqcup} {a.money}, _{\sqcup}b_{\sqcup}=_{\sqcup} {b.money}")
```

It works but... it can deadlock

Possible fix: fix order



- always acquire locks in the same order
- need a total order on locks
- we can use the getUniqueld (Example 1)

```
import SynchronizedProtectedUid.getUniqueId
class Account(val name: String, var money: Int) {
  val uid = getUniqueId()
}
```

Possible fix: fix order



- always acquire locks in the same order
- need a total order on locks
- we can use the getUniqueld (Example 1)

```
import SynchronizedProtectedUid.getUniqueId
class Account(val name: String, var money: Int) {
  val uid = getUniqueId()
}
```

Guarded blocks



Guarded block (for us)

a block of code that waits for a condition before running in a thread

Example 3: Thread pool with a queue of tasks

- Creating new threads in Java is expensive and avoidable
- Usually we re-use threads, by maintaining a set of waiting threads
- This set is call a thread pool
 - Scala already provides thread pools
 - We first create our own

```
import scala.collection.
object SynchronizedBadPool extends App {
// our set of tasks
 private val tasks = mutable.Queue[()=>Unit]()
// our single working thread
 val worker = new Thread {
   def poll(): Option[()=>Unit] =
     tasks.svnchronized {
      if (tasks.nonEmpty) Some(tasks.dequeue())
      else
                          None
  // keep on trying to run forever!
   override def run() = while (true)
     poll() match {
       case Some(task) => task()
       case None =>
   7
```

```
// starting the worker as
      a daemon
  worker.setName("Worker")
  worker, setDaemon(true)
  worker start()
  def asynchr(body: =>Unit) =
    tasks.synchronized {
      tasks.engueue(()=>body)
    3
  asynchr{ log("Hello") }
  asynchr{ log("_uworld!")}
  Thread.sleep(5000)
}
```



Daemon thread

- not the default
- have lower priority
- terminated automatically when JVM terminates
- in other words, do not prevent the JVM from terminating
- (the JVM terminates when 'normal' tasks terminate)



Busy-waiting is bad

- needlessly uses processor power (and drains the battery)
- after executing the previous code the worker will keep on running (unless you set in SBT set fork := true,)
- in general, we want the worker to enter a waiting state



synchronized + wait + notify

- these are methods that every Java/Scala object has
- wait:
 - needs the lock
 - puts the thread in a waiting state
 - releases the lock until activation
- notify:
 - needs the lock
 - activates all waiting threads



synchronized + wait + notify

- these are methods that every Java/Scala object has
- wait:
 - needs the lock
 - puts the thread in a waiting state
 - releases the lock until activation
- notify:
 - needs the lock
 - activates all waiting threads
- Note that the JVM can decide to call wait on its own spurious wakeups needing to re-enter the wait



```
object SynchronizedGuardedBlocks extends App {
 val lock = new AnyRef
 var message: Option[String] = None
 val greeter = thread {
   lock.synchronized {
     while (message == None) lock.wait() // non-busy waiting for a message
     log(message.get)
                                  // it will eventually log!
   }
 lock.synchronized {
   message = Some("Hello!")
   lock.notify()
                                    // awakes the (possibly) locked thread
 3
 greeter.join()
}
```



```
import scala.collection._
object SynchronizedPool extends App {
  private val tasks = mutable.Queue[()=>Unit]()
  object Worker extends Thread {
    setDaemon(true)
    def poll() = tasks.synchronized {
      while (tasks.isEmpty) tasks.wait()
                         // now using wait
      tasks.dequeue()
    3
    override def run() = while (true) {
      val task = poll()
      task()
    }
```

```
def asynchr(body: =>Unit) =
    tasks.svnchronized {
       tasks.enqueue(()=>body)
       // now notifying
       tasks.notify()
  asynchr{ log("Hello") }
  asynchr{ log("_uworld!")}
  Thread. sleep(500)
}
```

Worker.start()



- Our Worker can run forever (while-true)
- Terminates when the JVM terminates (daemon)
- Worker can be terminated earlier while waiting with Worker.interrupt()



- Our Worker can run forever (while-true)
- Terminates when the JVM terminates (daemon)
- Worker can be terminated earlier while waiting with Worker.interrupt()
- This triggers:
 - If it was waiting: an InterruptedException that can be handled
 - If it was not waiting: no exception and a flag Worker.isInterrupted becomes true



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- Terminates when the JVM terminates (daemon)
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- This triggers:
 - If it was waiting: an InterruptedException that can be handled
 - If it was not waiting: no exception and a flag Worker.isInterrupted becomes true
- Interrupts are needed if a thread does not awake with notify (e.g., it is doing blocking I/O)

Interrupting threads – alternative with graceful shutdown

```
object Worker extends Thread {
 var terminated = false
 // "manually" terminate when asked
 def poll(): Option[() => Unit] = tasks.synchronized {
    while (tasks.isEmpty && !terminated) tasks.wait()
   if (!terminated) Some(tasks.dequeue()) else None
 }
 import scala.annotation.tailrec
 Qtailrec override def run() = poll() match {
   case Some(task) => task(); run()
   case None =>
 }
 // "manually" ask to terminate
 def shutdown() = tasks.synchronized {
    terminated = true
   tasks.notify()
 3
```



FC

- using the @volatile annotation
- can be [atomically read] and [atomically modified]
- mostly used as status flag
- are never reordered in a thread
- writes are immediately visible to other threads
- very cheap to read
- not enough in many situations (e.g., getUniqueID)
- enough for previous example Slide 19

Example 4 – Batman



```
object Volatile extends App {
 class Page(val txt: String, var position: Int)
 val pages = for (i<- 1 to 5) yield
   new Page("Na" * (100 - 20 * i) + "__Batman!", -1)
 Qvolatile var found = false
 for (p <- pages) yield thread {</pre>
   var i = 0
   while (i < p.txt.length && !found) //</pre>
     if (p.txt(i) == '!') {
                               // Separate
       p.position = i
                                    // thread
       found = true
     } else i += 1
 3
 while (!found) {}
 log(s"results:__${pages.map(_.position)}")
```

The Java Memory Model overview



action α happens-before (HB) action β means action β sees the memory writes of action α

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- Transitivity: if α HB β and β HB $\gamma,$ then α HB γ
- Monitor locking: unlocking HB every subsequent locking (of the same lock)
- Volatile fields: writing to a volatile field HB every of its subsequent read

Data race

When a write to memory does not happen-before its intended read.

Immutable objects and final fields



```
class Foo( final val a: Int,
                  val b: Int.
                      c: Int)
// Encoding as Java:
class Foo {
  final private int a$;
  final private int b$;
  final private int c$;
  final public int a()
    { return a$; }
  public int b()
    { return b$; }
  public Foo(int a,
              int b.
             int c) {
    \{a\$ = a; b\$ = b; c\$ = c; \}
}
```

- Final fields: cannot be overridden
- val: cannot be updated
- vals are final
- Objects with only final fields
 - do not need synchronisation when shared (after constructed)
- Some collections are immutable (e.g. List), but contain non-final fields
 - need synchronisation when shared

Summary of operators



- Thread.sleep
- thr.start
- thr.join

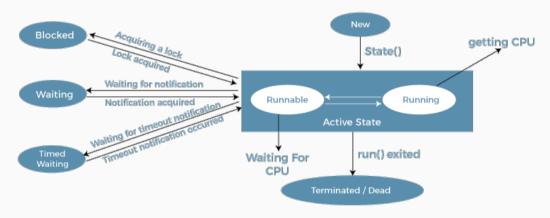
- thr.interrupt()
- thr.isInterrupted

@volatile var x

- lock.synchronized
- lock.wait
- lock.notify

Recall the Life-cycle of a thread





Life Cycle of a Thread

in https://static.javatpoint.com/core/images/life-cycle-of-a-thread.png

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The Java Memory Model overview